

# Passkeys With KeePassXC

An Introduction



### Outline

- What is KeePassXC?
- 2. The problem with passwords.
- 3. What are passkeys and how do they work?
- 4. Using passkeys with KeePassXC.









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图文

#### KeePassXC

Cross-platform Password Manager

Let KeePassXC safely store your passwords and auto-fill them into your favorite apps, so you can forget all about them.

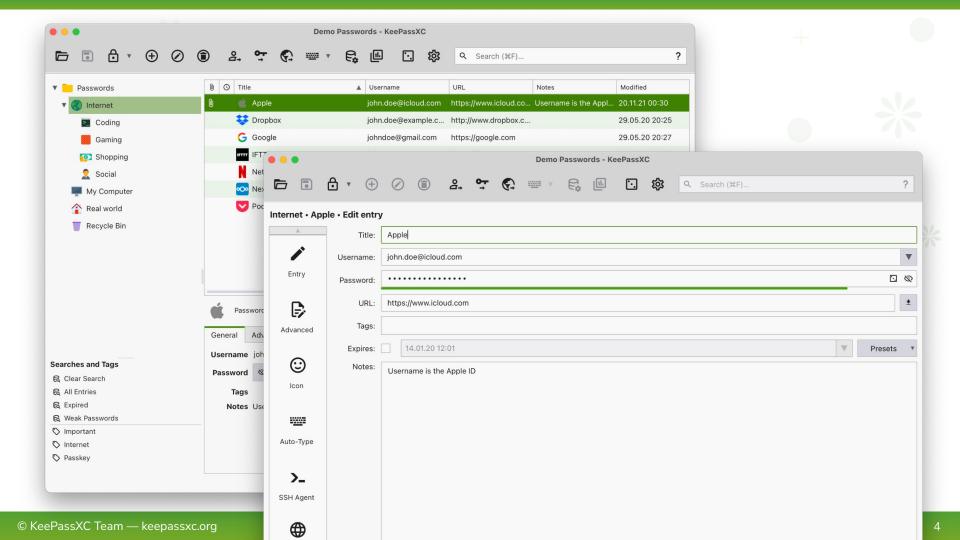
We do the heavy lifting in a no-nonsense, ad-free, tracker-free, and cloudfree manner. Free and open source.











## KeePass? KeePassX? KeePassXC?



2003–today keepass.info

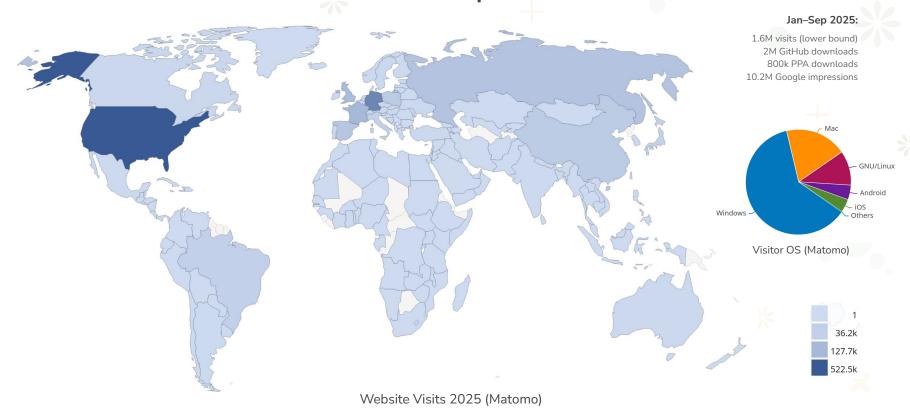


2005–2016 keepassx.org



2016–today keepassxc.org

# KeePassXC User World Map



# Password Security

#### Which password is best?

- qwerty
- asdf!123
- H3110 w0r1d!11
- 6F\$j5CA#@,OS
- opiCqjfWMwgSxeIQ8kNFovfup8YfwN
- ]÷O⁻¤Ú®mKá⁻û
- thaw unsmooth debate straggler syndrome tiara spotless poise

# Password Security

#### Which password is best?

- qwerty (~2 bit, « 1 second to crack)
- asdf!123 (~15 bit, 0.000001 seconds)
- H3110 w0r1d!11 (~32 bit, 0.1 seconds)
- **6F\$j5CA#@, OS** (~79 bit, 436.000 years)
- opiCqjfWMwgSxeIQ8kNFovfup8YfwN (~197 bit, 3.4x10<sup>23</sup> times age of the universe)
- ]÷O¯¤Ú®mKá¯û (~94 bit)
- thaw unsmooth debate straggler syndrome tiara spotless poise (~192 or 103 bit)

Cracking times: Single RTX 4090 with 22.000 MH/s (crude estimate)

# Password Security

#### Which password is best?

- qwerty (~2 bit, « 1 second to crack)
- asdf!123 (~15 bit, 0.000001 seconds)
  H3110 w0r1d!11 (~32 bit, 0.2 deconds)
  6F\$j5CA#@, OS (~70 bit, 410900 years)
- opiCqjfWMwgSxeIQ8lQYFovfup8YfwN (~162 bit, 4.7x10<sup>14</sup> times age of the universe)
- ]÷O¯¤Ú®mKá¯û (30°0 bit)
- thaw unsmooth debate straggler syndrome tiara spotless poise (~192 or 109 bit)

Cracking times: Single RTX 4090 with 22.000 MH/s (crude estimate)

#### Problems With Passwords

- Hard to remember (PMs solve this).
- Cumbersome to use (PMs try to solve this).
- Prone to phishing (PM browser extensions make phishing harder).
- Reuse makes users vulnerable (PMs discourage reuse).
- Must be changed when service compromised (PMs cannot solve this).

→ Passkeys are meant to solve all of the above.

# Passkeys Aren't Perfect

- Web service must support it.
- Cannot replace your banking PIN, locker combination, etc.
- Not meant for offline data encryption (but can be done).
- Hardware authenticators have limited storage capacity.
- Software authenticators vulnerable to malware.
- Import / export formats still work in progress.
- Standard allows enforcement of specific authenticators (vendor lock-in).

## Passkeys: A Short History

- 2013: FIDO Alliance founded
- 2014: Universal 2nd Factor (U2F) 1.0 standard released
- 2015: FIDO 2.0 proposal submitted to W3C
- 2017: U2F 1.2 standard released



Image credits: Yubico

- 2018: Client To Authenticator Protocol 2.0 (CTAP2) standard proposed
   U2F renamed to CTAP1
  - 2019: W3C WebAuthn Level 1 recommendation published
- 2020: Most major browsers support FIDO2
- 2022: Chrome and Safari ship native Passkey support
- 2024: Firefox supports native Passkeys as last major browser

# FIDO? CTAP? WebAuthn? Passkey?

#### FIDO (Fast IDentity Online)

- FIDO Alliance = The Consortium.
- FIDO1 + FIDO2 = The parent specifications for everything below.

#### CTAP

- Protocol for talking to FIDO authenticators.
- o CTAP1 is the same as U2F.
- CTAP2 is one of two parts of the FIDO2 specification.

#### WebAuthn

- W3C-standardised API for initiating passwordless user authentication (via CTAP).
- WebAuthn is the second part of the FIDO2 specification.

#### Passkeys

- Marketing term by Apple without clear definition.
- Usually refers to (discoverable) FIDO2 credential key pair.
- Windows, macOS, iOS, Android have native passkey support (Linux in the works).



Principles of secure information systems (there are two more):

- Confidentiality: The message must remain secret
- Integrity: The message must be unchanged
- Authenticity: The source must be trustworthy

Let P be the set of all texts (plain texts), K the set of all keys, C the set of all encrypted texts (cipher texts), and  $e_k$ ,  $d_k$  two functions:

$$\begin{array}{ll} e_k: P \longrightarrow C \\ \\ d_k: C \longrightarrow P \end{array} \qquad \text{with} \qquad d_k(e_k(x)) = x, \quad x \in P, \ k \in K \end{array}$$

**Problem:** Transmission of the (secret) key k.

Idea: Alice and Bob each have two keys  $k_{pub}$  (public) and  $k_{priv}$  (private) so that

$$d_k \left( e_k \left( x \right) \right) = e_k \left( d_k \left( x \right) \right)$$

#### Steps of asymmetric encryption:

- 1. Alice and Bob choose keys  $k_{pol}^A$ ,  $k_{priv}^A$  and  $k_{pol}^B$ ,  $k_{priv}^B$ .
- 2. Both publish their public keys  $k_{pub}^A$ ,  $k_{pub}^B$ .
- 3. Alice sends message x as  $y = e_{k^B}(x)$  to Bob.
- 4. Bob decrypts y and gets  $x = d_{k_{\underline{B}}}(y)$ .

Source: Stein, Bevendorff – webis.de (German)

How does Alice know that Bob and not Eve is the sender of x?

#### Digital signatures:

Let  $h: P \to N$  be a hash function, which calculates for x a unique characterisation h(x) of fixed length ("message digest").

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#### Digital signatures:

Let  $h: P \to N$  be a hash function, which calculates for x a unique characterisation h(x) of fixed length ("message digest").

- Alice calculates for x the hash value h(x).
- Alice encrypts h(x) as  $y_h = e_{k^A}(h(x))$ .
- Alice sends  $y = e_{k_{pol}^B}(x+y_h)$  to Bob. Bob decrypts  $x+y_h = d_{k_p^B}(y)$ .
- Bob calculates h(x) and  $d_{k^A}(y_h)$  and compares values.

Source: Stein, Bevendorff - webis.de (German)

# Universal Second Factor (U2F)

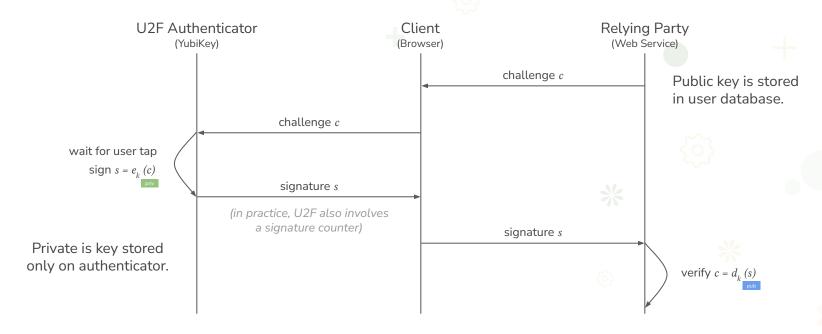
Traditionally, authentication factors are defined as:

- Something you know (password, PIN, ...)
- Something you have (smart card, OTP generator, TAN list, ...)
- Something you are (facial recognition, fingerprint, other biometrics, ...)

→ Ideally, you have at least two.

# Universal Second Factor (U2F)

U2F devices model "something you have" with public key crypto:

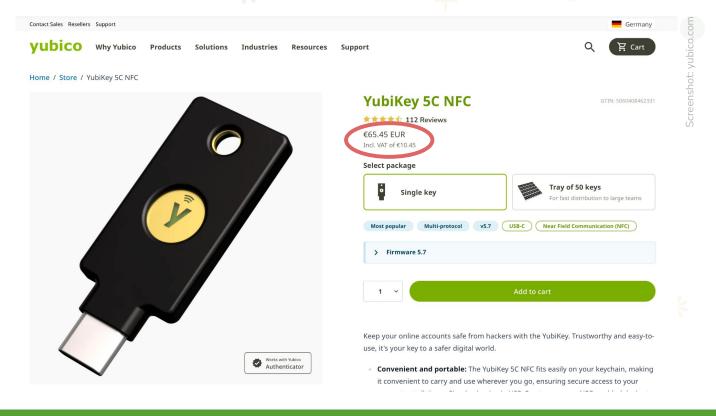


# Going Passwordless With FIDO2

#### Main improvements over FIDO U2F:

- Proper API standardisation (WebAuthn)
- Individual (pass)keys per user and service
- Discoverable credentials (resident keys)
- Passwordless MFA via PIN or biometrics
- Platform authenticators → KeePassXC

# Going Passwordless With FIDO2



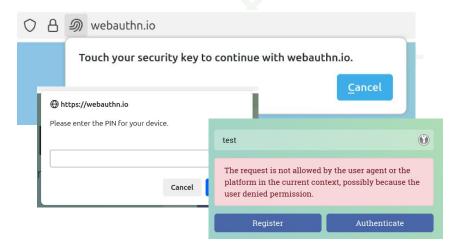
# Passkey Authentication Workflow



Authenticate

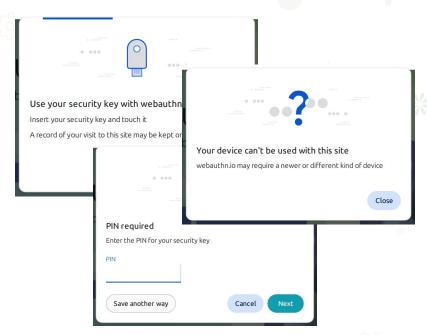
# Linux: No Native Passkeys (Yet)

#### Firefox



YubiKey: V Platform: X

#### Chrome



# Passkeys With KeePassXC

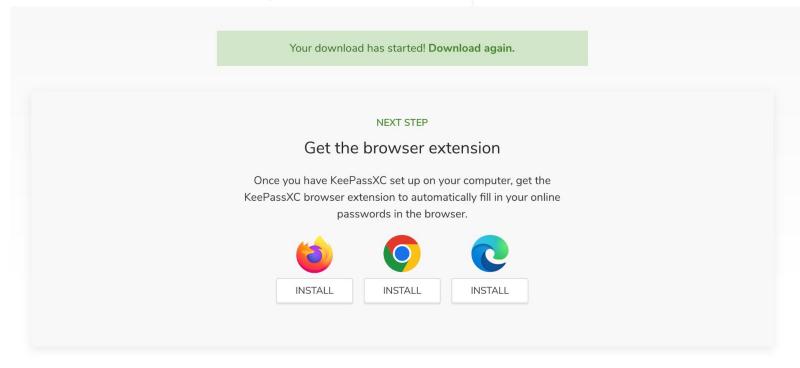
keepassxc.org/download/



#### Linux Tips:

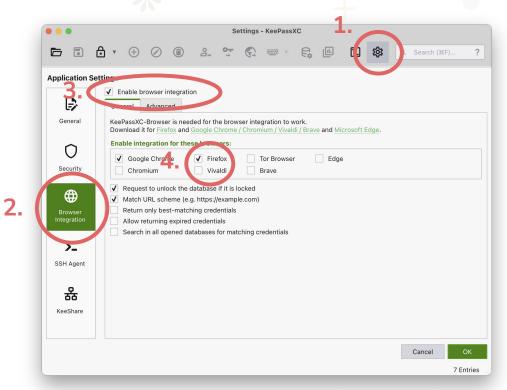
- Flatpak or PPA are fine
- Applmage if you need to
- Distro package if not too old
- Avoid Snaps

# Passkeys With KeePassXC

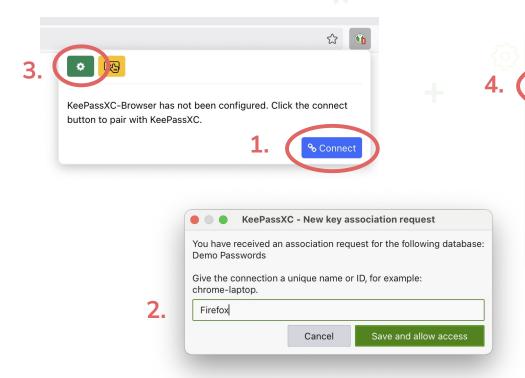


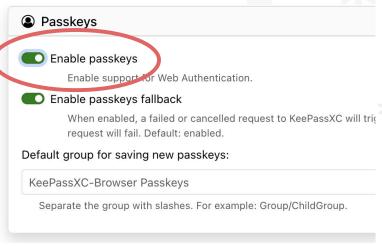
**Important:** Snap / Flatpak browsers are not supported!

# **Enable Browser Integration**



# Connect KeePassXC-Browser





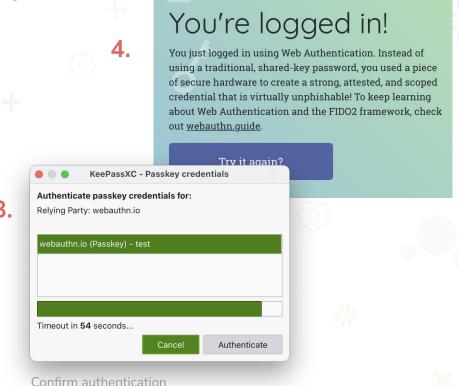
# Register New Passkey



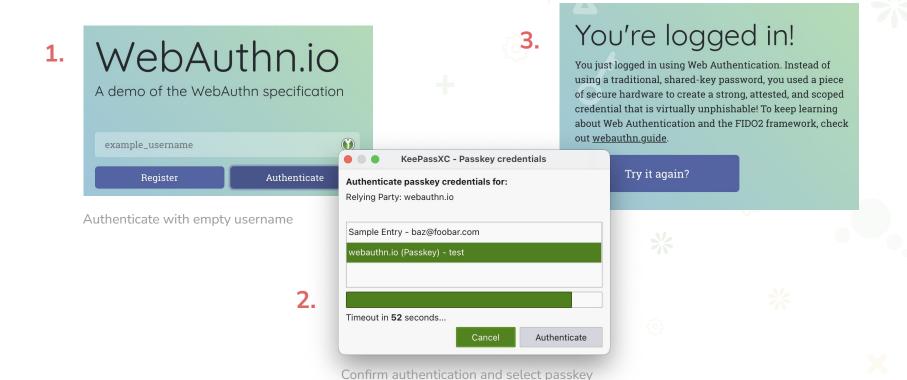
# Authenticate Via Passkey

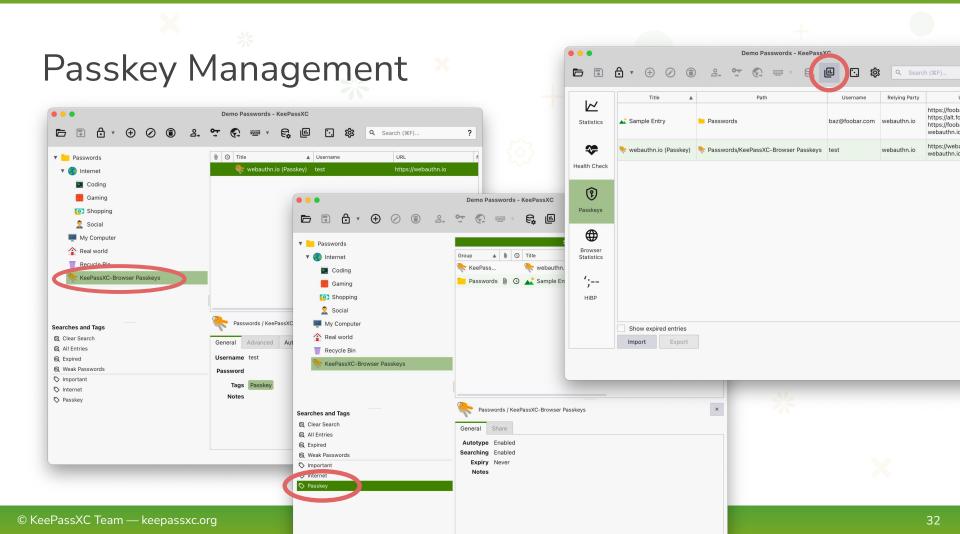


Select passkey from menu

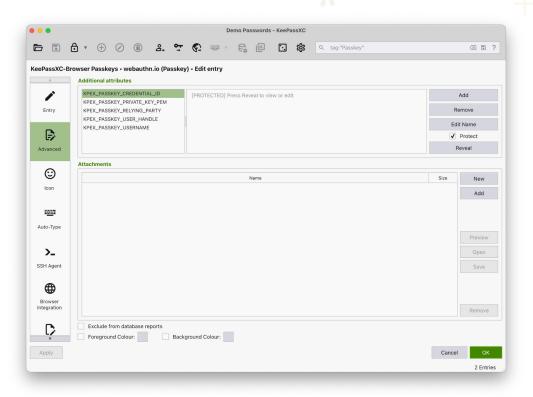


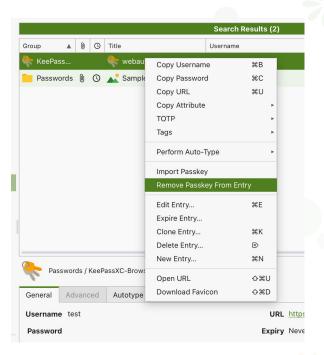
# Discoverable Credentials





# Passkey Management





# Summary

Passkeys with KeePassXC...

- ... allow passwordless login (except for the master password 😉).
- ... prevent credential phishing.
- ... prevent password leakage.
- ... are portable and easy to back up.

#### Thanks!

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